# Universal Conditional Gradient Sliding 

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## Problem Setting

## Convex Optimization

Our problem of interest is computing an $\varepsilon$-solution $\tilde{x} \in X$ to

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\begin{equation*}
f^{*}:=\min _{x \in X} f(x) \tag{CO}
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such that $f(\tilde{x})-f^{*}<\varepsilon$.

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such that $f(\tilde{x})-f^{*}<\varepsilon$.
Here,

- $f$ is a real-valued, convex function
- $\mathbb{R}^{n}$ is a high dimensional space
- $X \subseteq \mathbb{R}^{n}$ closed, bounded, and convex


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- differentiability of $f$
- smoothness of $f$


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- Lipschitz continuity of $f$
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- smoothness of $f$
- Lipschitz continuity of $f$
- strong convexity of $f$

For now, let us assume $\nabla f$ exists and is Lipschitz continuous with Lipschitz constant $L$, i.e. $f$ is L-smooth, and that a projection onto $X$ is computationally feasible.

## Solving (CO)

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How fast is fast and how do we measure this?
$\rightarrow$ count the number of expensive operations
Example: Gradient Descent (GD)

$$
\begin{aligned}
x_{k} & =\underset{u \in X}{\operatorname{argmin}}\left\|u-\left(x_{k-1}-\frac{1}{\eta_{k}} \nabla f\left(x_{k-1}\right)\right)\right\|^{2} \\
& =\underset{u \in X}{\operatorname{argmin}}\left\langle\nabla f\left(x_{k-1}\right), u\right\rangle+\frac{\eta_{k}}{2}\left\|u-x_{k-1}\right\|^{2}
\end{aligned}
$$

- for properly chosen $\eta_{k}$, GD achieves an $\varepsilon$-solution in $\mathcal{O}(1 / \varepsilon)$ iterations
- only expensive operation is gradient evaluation, and GD uses 1 per iteration


## Solving (CO)

## Algorithm 1 Nesterov's accelerated gradient descent (NAGD)

Start: Choose $x_{0} \in X$. Set $y_{0}:=x_{0}$
for $k=1, \ldots, N$ do

$$
\begin{aligned}
& z_{k}=\left(1-\gamma_{k}\right) y_{k-1}+\gamma_{k} x_{k-1}, \\
& x_{k}=\underset{u \in X}{\operatorname{argmin}}\left\langle\nabla f\left(z_{k}\right), u\right\rangle+\frac{\eta_{k}}{2}\left\|u-x_{k-1}\right\|^{2}, \\
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Output $y_{N}$.

- minimizes linear approximation proximal problem
- subproblem is still a projection
- reduces to gradient descent when $\gamma_{k} \equiv 1$


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## end for

Output $y_{N}$.

- minimizes linear approximation proximal problem
- subproblem is still a projection
- reduces to gradient descent when $\gamma_{k} \equiv 1$
- computes $\varepsilon$-solution in only $\mathcal{O}(1 / \sqrt{\varepsilon})$ iterations
- requires knowledge of $L$ to set $\eta_{k}$ appropriately
- is optimal for solving problems such as (CO) under first order oracle ([1])


## Solving (CO)

But what if the projection is not so easy?

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But what if the projection is not so easy?

- certain sets can be as difficult to project to as the underlying problem is to solve
$>X=\operatorname{conv}\left(v_{1}, \ldots, v_{p}\right)$
$>X=\left\{Y \in \mathbb{R}^{n \times n}: \operatorname{tr}(Y)=1, Y \succeq 0\right\}$
- NAGD is of no use when projection is more difficult than (CO)
- want to design algorithms that do not require difficult optimizations over $X$, i.e. projection free methods


## Projection Free Methods for Solving (CO)

## Algorithm 2 Conditional Gradient (CG)

Start: Choose $y_{0} \in X$.
for $k=1, \ldots, N$ do

$$
\begin{aligned}
& x_{k}=\underset{x \in X}{\operatorname{argmin}}\left\langle\nabla f\left(y_{k-1}\right), x\right\rangle \\
& y_{k}=\left(1-\alpha_{k}\right) y_{k-1}+\alpha_{k} x_{k}
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end for
Output $y_{N}$.

- solves a linear optimization (LO) rather than a projection
$>$ when $X$ is convex hull, the LO is a linear program
$>$ when $X$ is standard spectrahedron, the LO is a smallest eigenvalue problem
- requires $\mathcal{O}(1 / \varepsilon)$ number of iterations to obtain $\varepsilon$-solution [2]
- more gradient evaluations and the addition of linear optimizations, but no projections at all
- optimal number of linear optimizations


## Projection Free Methods for Solving (CO)

## Question

Comparing CG to NAGD, we increase in the complexity of gradient evaluations necessary. Is it possible to keep the gradient evaluations unchanged while being projection free?

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Comparing CG to NAGD, we increase in the complexity of gradient evaluations necessary. Is it possible to keep the gradient evaluations unchanged while being projection free?

The answer is yes! Simply solve $x_{k}$ subproblem with a projection free algorithm.


## Projection Free Methods for Solving (CO)

$$
\begin{aligned}
& \text { Algorithm } 3 \text { Conditional Gradient Sliding (CGS) } \\
& \hline \text { Start: Choose } x_{0} \in X . \text { Set } y_{0}:=x_{0} \\
& \text { for } k=1, \ldots, N \text { do } \\
& \qquad \begin{aligned}
z_{k} & =\left(1-\gamma_{k}\right) y_{k-1}+\gamma_{k} x_{k-1}, \\
x_{k} & =C G\left(\nabla f\left(\underline{x}_{k}\right), x_{k-1}, \eta_{k}, \varepsilon_{k}\right) \\
y_{k} & =\left(1-\gamma_{k}\right) y_{k-1}+\gamma_{k} x_{k} .
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\end{aligned}
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$$

end for
Output $y_{N}$.

- solves $x_{k}$ subproblem approximately with linear optimizations only
- if parameters are chosen properly, CGS computes $\varepsilon$-solution in $\mathcal{O}(1 / \sqrt{\varepsilon})$ gradient evaluations and $\mathcal{O}(1 / \varepsilon)$ linear optimizations [3]
- requires $L$ for setting of $\eta_{k}$


## Smoothness

A key feature of all the above algorithms is the assumption that the gradient of $f$ is Lipschitz with constant L, i.e,

$$
f(x) \leq f(u)+\langle\nabla f(u), x-u\rangle+\frac{L}{2}\|x-u\|^{2} .
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$$

This combined with convexity can be leveraged to design efficient optimization methods. However, we may not always have such luxury.

- $f(x)=\lambda\|x\|$
- $f(x)=\max _{y \in \Delta_{m}}\langle x, A y\rangle$


## Smoothness Relaxation

## Relaxed Assumption - Hölder Smooth

Assume that there exists a Hölder exponent $\nu \in[0,1]$ and constant $M_{\nu}>0$ such that

$$
f(y) \leq f(x)+\langle\nabla f(x), y-x\rangle+\frac{M_{\nu}}{1+\nu}\|x-y\|^{1+\nu}, \forall x, y \in X
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$$

This is a generalization of Lipschitz continuous gradient. In particular,

- any convex smooth $f$ with Lipschitz continuous gradient $M_{1}$ is Hölder smooth with $\nu=1$
- any convex nonsmooth Lipschitz continuous $f$ with is Hölder smooth with $\nu=0$
- any convex smooth $f$ satisfying

$$
\|\nabla f(y)-\nabla f(x)\| \leq M_{\nu}\|y-x\|^{\nu}, \forall x, y \in X
$$

is Hölder smooth with $\nu \in(0,1)$

## Solving (CO) with Sliding

## Algorithm 4 Fast Gradient Method (FGM)

$$
\begin{aligned}
& \text { Start: Choose } x_{0} \in X \text { and } \varepsilon>0 \text {. Set } y_{0}=x_{0} \\
& \text { for } k=1, \ldots, N \text { do } \\
& \text { Decide } L_{k}>0 \text { satisfying } \\
& \qquad f\left(y_{k}\right) \leq f\left(z_{k}\right)+\left\langle\nabla f\left(z_{k}\right), y_{k}-z_{k}\right\rangle+\frac{L_{k}}{2}\left\|y_{k}-z_{k}\right\|^{2}+\frac{\varepsilon}{2} \gamma_{k}
\end{aligned}
$$

where

$$
\begin{aligned}
& z_{k}=\left(1-\gamma_{k}\right) y_{k-1}+\gamma_{k} x_{k-1}, \\
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\end{aligned}
$$

end for
Output $y_{N}$.

- achieves $\varepsilon$-solution in $\mathcal{O}\left((1 / \varepsilon)^{\frac{2}{1+3 \nu}}\right)$ iterations
- also optimal


## Comparison of Algorithms

We repeat the improvements

- rather than minimizing $f$ with FGM, we can apply CG to minimize using linear optimizations instead of projections
- CG for functions with Hölder continuous gradients requires $\mathcal{O}\left((1 / \varepsilon)^{\nu}\right)$ iterations [4]
- we can preserve the $\mathcal{O}\left((1 / \varepsilon)^{\frac{2}{1+3 \nu}}\right)$ gradient evaluations by approximately solving the $x_{k}$ subproblem in FGM using CG


## Numerical Example

## Algorithm 5 Universal Conditional Gradient Sliding (UCGS)

## Start: Choose $x_{0} \in X$ and $\varepsilon>0$. Set $y_{0}=x_{0}$

for $k=1, \ldots, N$ do
Decide $L_{k}>0$ satisfying

$$
f\left(y_{k}\right) \leq f\left(z_{k}\right)+\left\langle\nabla f\left(z_{k}\right), y_{k}-z_{k}\right\rangle+\frac{L_{k}}{2}\left\|y_{k}-z_{k}\right\|^{2}+\frac{\varepsilon}{2} \gamma_{k}
$$

where

$$
\begin{aligned}
& z_{k}=\left(1-\gamma_{k}\right) y_{k-1}+\gamma_{k} x_{k-1}, \\
& x_{k}=A C G\left(\nabla f\left(z_{k}\right), x_{k-1}, \eta_{k}, \varepsilon_{k}, \delta_{k}\right) \\
& y_{k}=\left(1-\gamma_{k}\right) y_{k-1}+\gamma_{k} x_{k} .
\end{aligned}
$$

Terminate if

$$
\max _{x \in X} f\left(y_{k}\right)-\ell_{k}(x) \leq \varepsilon
$$

where

$$
\ell_{k}(x)=\Gamma_{k} \sum_{i=1}^{k} \frac{\gamma_{i}}{\Gamma_{i}}\left(f\left(z_{i}\right)+\left\langle\nabla f\left(z_{i}, x-z_{i}\right\rangle\right)\right.
$$

end for
Output $y_{N}$.

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## Properties of UCGS

- contains a stopping condition


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- maintains the $\mathcal{O}\left((1 / \varepsilon)^{\frac{2}{1+3 \nu}}\right.$ ) gradient evaluations established in FGM for an $\varepsilon$-solution


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- maintains the $\mathcal{O}\left((1 / \varepsilon)^{\frac{2}{1+3 \nu}}\right.$ ) gradient evaluations established in FGM for an $\varepsilon$-solution
- improves the number of linear optimizations required of CG to $\mathcal{O}\left((1 / \varepsilon)^{\frac{4}{1+3 \nu}}\right)$


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- improves the number of linear optimizations required of CG to $\mathcal{O}\left((1 / \varepsilon)^{\frac{4}{1+3 \nu}}\right)$
- allows linear optimization problems to be solved approximately


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- maintains the $\mathcal{O}\left((1 / \varepsilon)^{\frac{2}{1+3 \nu}}\right.$ ) gradient evaluations established in FGM for an $\varepsilon$-solution
- improves the number of linear optimizations required of CG to $\mathcal{O}\left((1 / \varepsilon)^{\frac{4}{1+3 \nu}}\right)$
- allows linear optimization problems to be solved approximately
- achievable by novel parameter choice


## Summary of Advantages

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Advantages over CGS

- no longer requires knowledge of $L$ in the smooth case
- provides additional application for $\nu \in[0,1)$
- allows usage of inexact linear optimization solvers
- allows for possibility of early termination with exit criterion


## Numerical Experiments - Convex Hull

We consider the problem

$$
\min _{x \in \operatorname{conv}(V)} f(x):=\|A x-b\|_{2}
$$

with $V=\left\{v_{1}, \ldots, v_{p}\right\} \subseteq \mathbb{R}^{n}, \operatorname{conv}(V):=\left\{x \in \mathbb{R}^{n}: \exists \lambda \in \Delta_{p}\right.$ s.t. $\left.x=\sum_{j=1}^{p} \lambda_{i} v_{i}\right\}$, and $\Delta_{p}:=\left\{\lambda \in \mathbb{R}^{p}: \sum_{i=1}^{p} \lambda_{i}=1, \lambda_{i} \geq 0\right\}$ is the standard simplex.

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|  |  | UCGS |  |  |  | CG |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $n$ | $d$ | GE | LO | Time | Error | Iter | Time | Error |
| 2500 | 0.2 | 66 | 2690 | 6.71 | $9.945 e-4$ | 572 | 13.42 | $9.7086 e 1$ |
| 2500 | 0.4 | 60 | 3679 | 9.08 | $9.976 e-4$ | 524 | 18.17 | $1.404 e 2$ |
| 2500 | 0.6 | 62 | 245 | 2.64 | $9.678 e-4$ | 146 | 5.29 | $5.598 e 2$ |
| 2500 | 0.8 | 57 | 3176 | 8.45 | $9.768 e-4$ | 399 | 16.93 | $2.400 e 2$ |
| 5000 | 0.2 | 71 | 286 | 7.13 | $9.882 e-4$ | 178 | 14.32 | $6.037 e 2$ |
| 5000 | 0.4 | 42 | 52 | 4.89 | $9.585 e-4$ | 84 | 9.81 | $1.689 e 3$ |
| 5000 | 0.6 | 68 | 4564 | 36.14 | $9.727 e-4$ | 483 | 72.40 | $3.527 e 2$ |
| 5000 | 0.8 | 67 | 419 | 12.91 | $9.815 e-4$ | 161 | 25.94 | $1.165 e 3$ |
| 10000 | 0.2 | 85 | 12269 | 150.51 | $9.96 e-4$ | 915 | 301.21 | $2.449 e 2$ |
| 10000 | 0.4 | 69 | 12614 | 157.39 | $9.916 e-4$ | 636 | 315.27 | $4.734 e 2$ |
| 10000 | 0.6 | 70 | 16063 | 205.87 | $9.821 e-4$ | 653 | 412.14 | $5.423 e 2$ |
| 10000 | 0.8 | 69 | 12707 | 180.65 | $9.862 e-4$ | 473 | 361.73 | $8.162 e 2$ |

## Numerical Experiments - Spectrahedron

For our second experiment, we solve the problem

$$
\min _{X \in \mathrm{Spe}_{n}} f(X):=\sum_{i=1}^{m}\left\|X-A_{i}\right\|_{2}
$$

where Spe $_{n}:=\left\{X \in \mathbb{R}^{n \times n}: \operatorname{tr}(X)=1, X \succeq 0\right\}$ and $A_{i} \in$ Spe $_{n}$ for each $i=1, \ldots, m$.

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|  |  | UCGS |  |  |  | CG |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $n$ | $m$ | GE | LO | Time | Error | Iter | Time | Error |
| 50 | 50 | 1354 | 8493 | 9.87 | $9.992 e-4$ | 6908 | 19.74 | $6.073 e-3$ |
| 50 | 100 | 1767 | 11138 | 13.09 | $9.994 e-4$ | 7038 | 26.19 | $1.172 e-2$ |
| 50 | 200 | 2425 | 15173 | 25.39 | $9.995 e-4$ | 8273 | 50.79 | $2.271 e-2$ |
| 100 | 50 | 1836 | 13056 | 159.61 | $9.980 e-4$ | 11648 | 319.25 | $3.225 e-3$ |
| 100 | 100 | 2347 | 16816 | 216.59 | $9.990 e-4$ | 13372 | 433.20 | $5.634 e-3$ |
| 100 | 200 | 3296 | 23836 | 310.16 | $9.984 e-4$ | 16053 | 620.36 | $9.892 e-3$ |
| 200 | 50 | 1722 | 33673 | 470.71 | $9.989 e-4$ | 15966 | 941.43 | $3.308 e-3$ |
| 200 | 100 | 2314 | 46323 | 730.69 | $9.994 e-4$ | 17033 | 1461.42 | $6.870 e-3$ |
| 200 | 200 | 3154 | 64511 | 1086.42 | $9.992 e-4$ | 19762 | 2172.85 | $1.015 e-2$ |

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